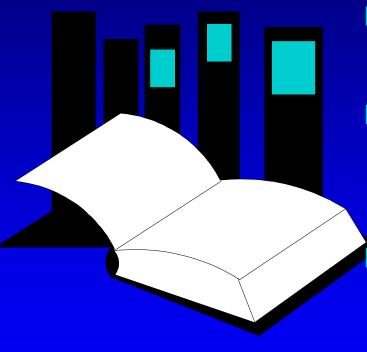
CSC212 Data Structure



Lecture 14 Binary Search Trees

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City College of New York

Binary Search Trees



- □ One of the tree applications in Chapter 10 is binary search trees.
- ☐ In Chapter 10, binary search trees are used to implement bags and sets.
 - This presentation illustrates how another data type called a **dictionary** is implemented with binary search trees.

Binary Search Tree Definition

- □ In a binary search tree, the entries of the nodes can be compared with a strict weak ordering. Two rules are followed for every node n:
 - ☐ The entry in node n is NEVER less than an entry in its left subtree
 - ☐ The entry in the node n is less than every entry in its right subtree.

- □ A dictionary is a collection of items, similar to a bag.
- But unlike a bag, each item has a string attached to it, called the item's <u>key</u>.

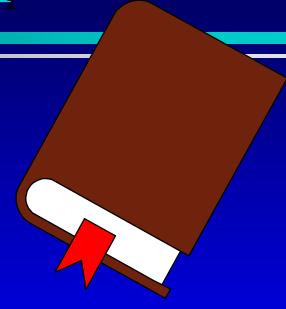


- □ A dictionary is a collection of items, similar to a bag.
- But unlike a bag, each item has a string attached to it, called the item's key.

Example:

The items I am storing are records containing data about a state.





- □ A dictionary is a collection of items, similar to a bag.
- But unlike a bag, each item has a string attached to it, called the item's key.

Example:

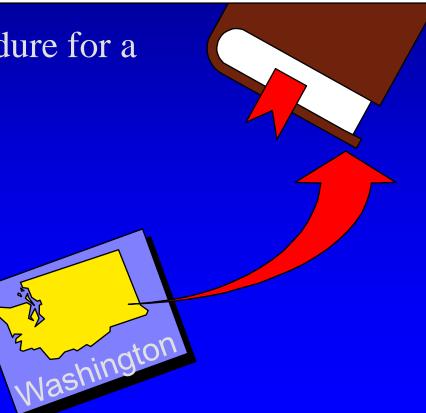
The **key** for each record is the name of the state.



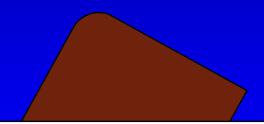


void Dictionary::insert(The key for the new item, The new item);

The insertion procedure for a dictionary has two parameters.



□ When you want to retrieve an item, you specify the key...

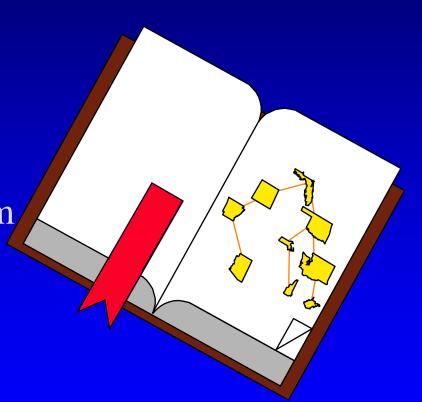


Item Dictionary::retrieve("Washington");



□ When you want to retrieve an item, you specify the key... ... and the retrieval procedure returns the item. Item

We'll look at how a binary tree can be used as the internal storage mechanism for the dictionary.

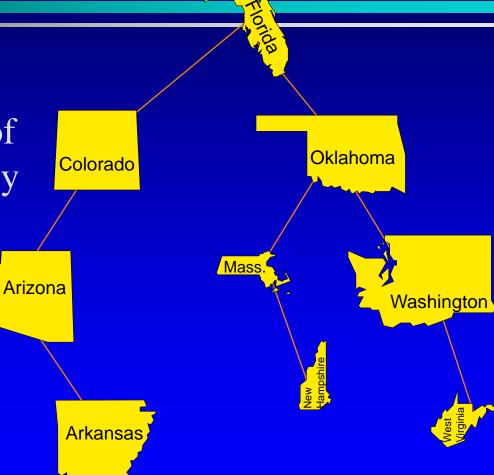


The data in the dictionary will be stored in a binary tree, with each node containing an item and a key.



Storage rules:

Every key to the **left** of a node is alphabetically **before** the key of the node.



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Every key to the **left** of a node is alphabetically **before** the key of the node.

Example:

- ' Massachusetts' and
- ' New Hampshire' are alphabetically before 'Oklahoma'

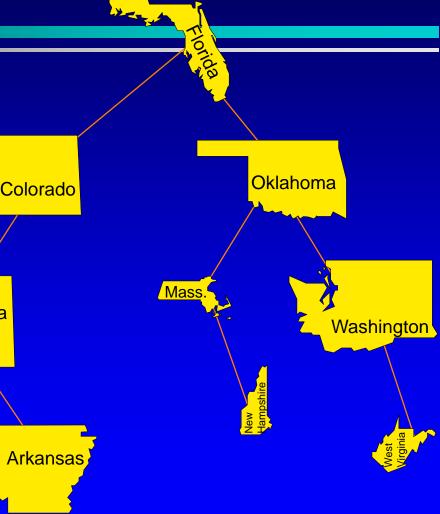


Arizona

Storage rules:

Every key to the <u>left</u> of a node is alphabetically <u>before</u> the key of the node.

Every key to the right of a node is alphabetically after the key of the node.



Arizona

Storage rules:

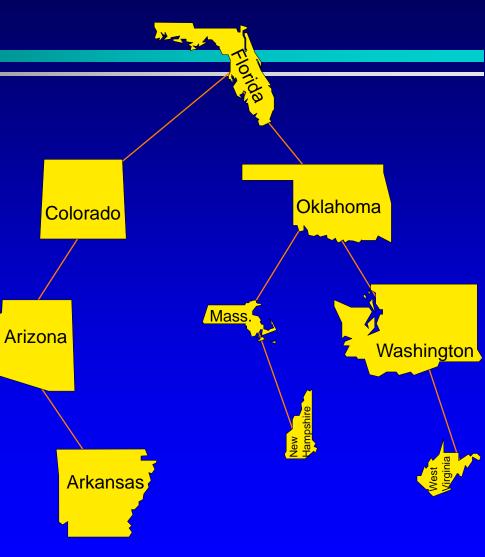
Every key to the **left** of a node is alphabetically **before** the key of the node.

☐ Every key to the right of a node is alphabetically after the key of the node.



Retrieving Data

- ☐ If the current node has the key, then stop and retrieve the data.
- ☐ If the current node's key is too large, move left and repeat 1-3.
- ☐ If the current node's key is too small, move right and repeat 1-3.



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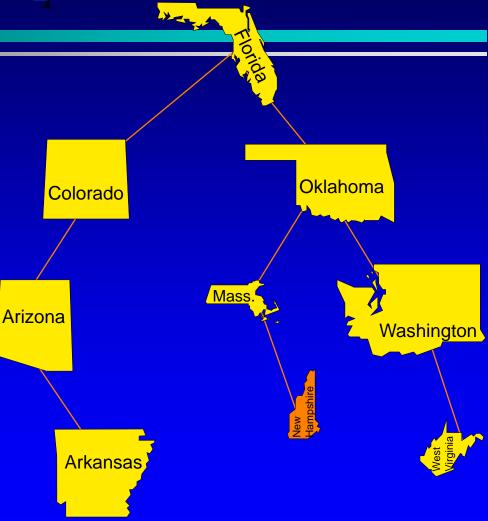
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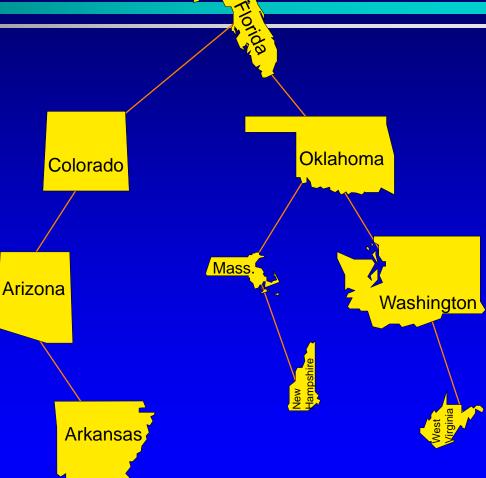
- ☐ If the current node has the key, then stop and retrieve the data.
- If the current node's key is too <u>large</u>, move <u>left</u> and repeat 1-3.
- ☐ If the current node's key is too small, move right and repeat 1-3.



Adding a New Item with a Given Key

☐ Pretend that you are trying to find the key, but stop when there is no node to move to.

☐ Add the new node at the spot where you would have moved to if there had been a node.



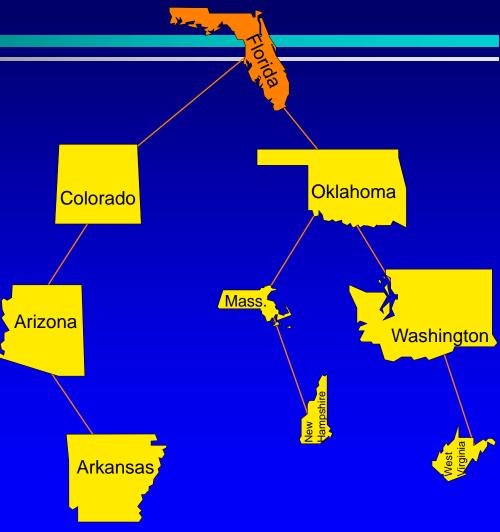


- Pretend that you are trying to find the key, but stop when there is no node to move to.
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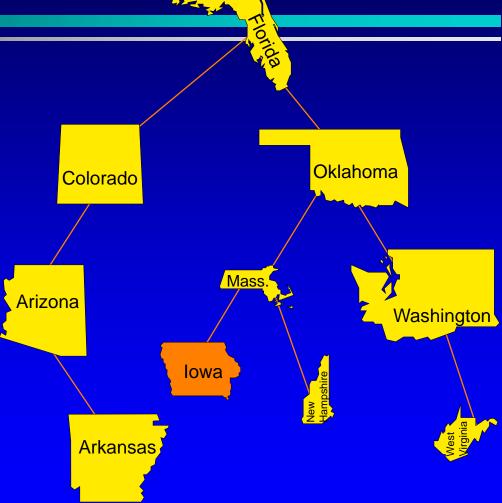




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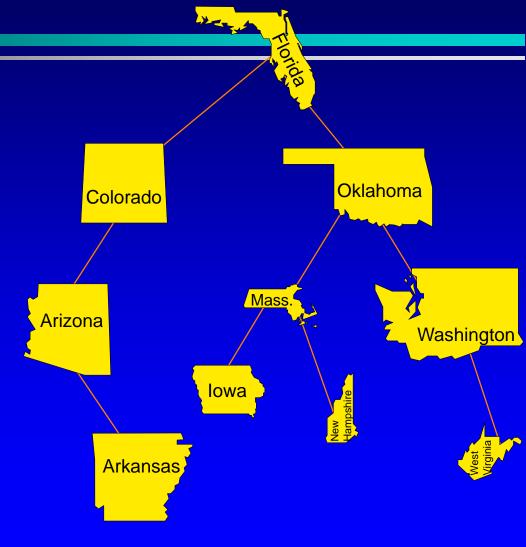


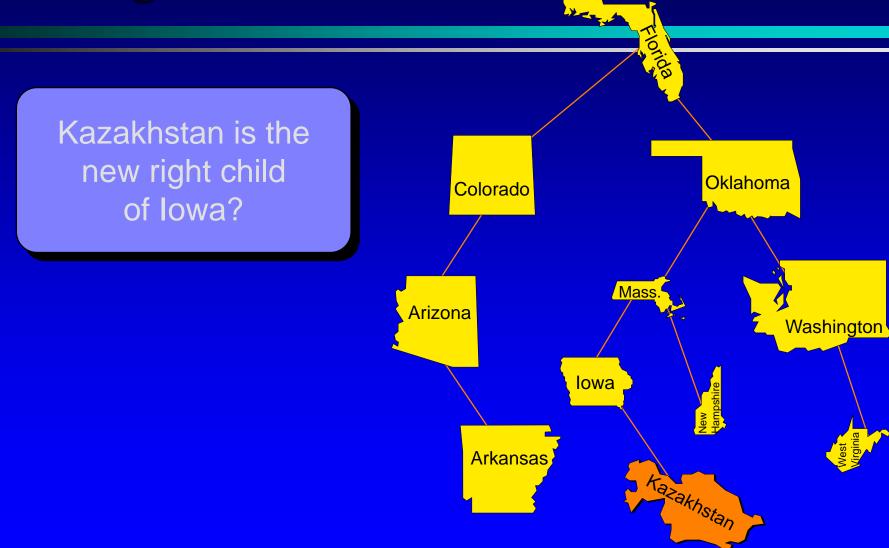
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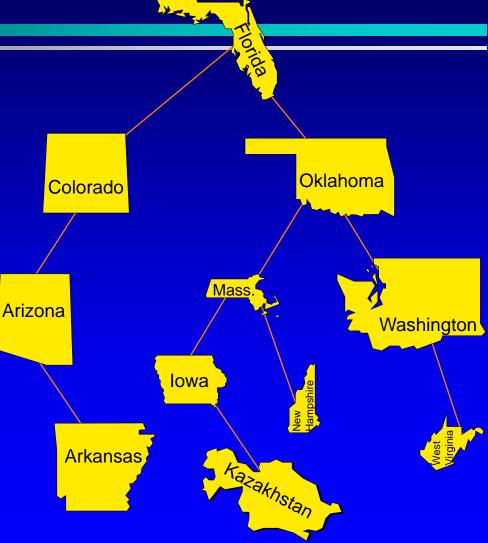
Where would you add this state?



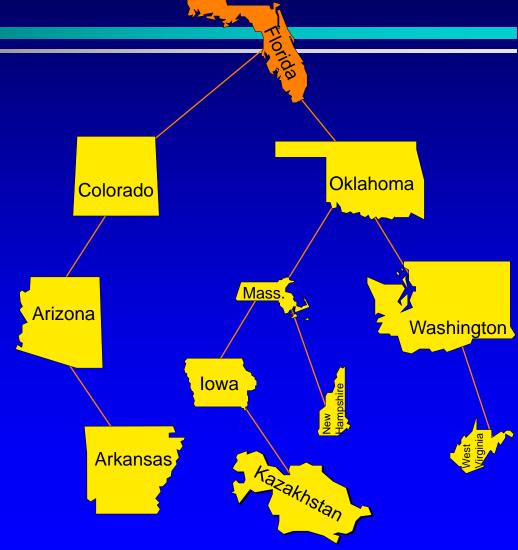


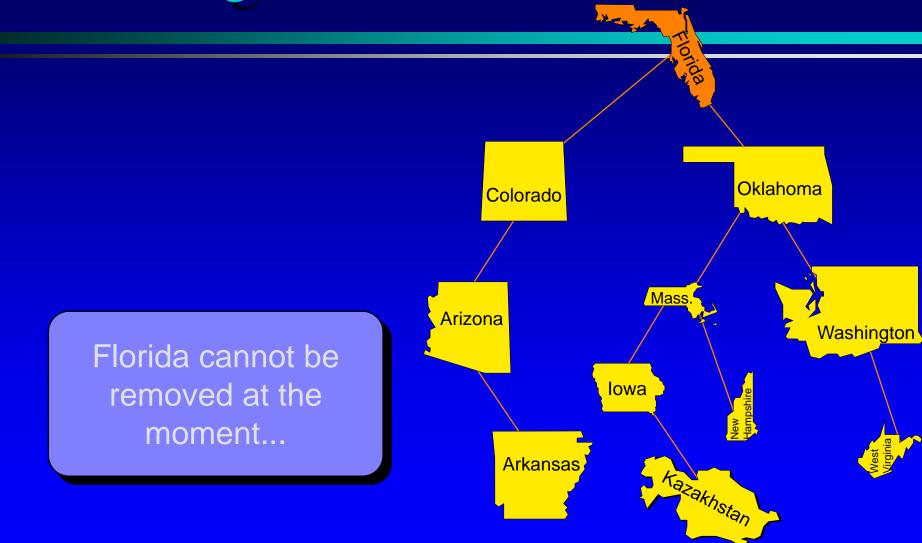
Removing an Item with a Given Key

- ☐ Find the item.
- ☐ If necessary, swap the item with one that is easier to remove.
- ☐ Remove the item.



Find the item.



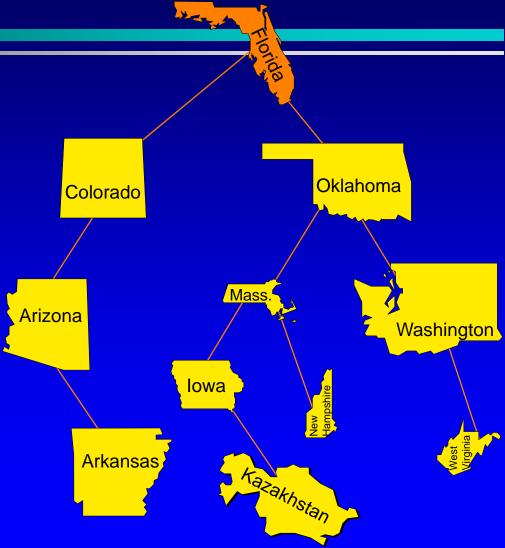


... because removing
Florida would
break the tree into
two pieces.



☐ If necessary, do some rearranging.

The problem of breaking the tree happens because Florida has 2 children.



☐ If necessary, do some rearranging.

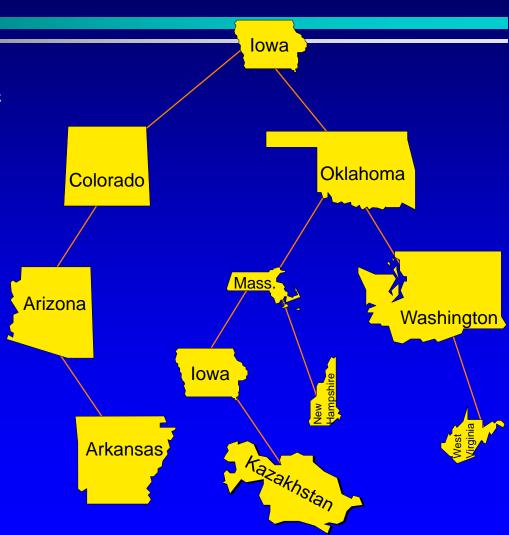
For the rearranging, take the **smallest** item in the right subtree...

Oklahoma Colorado Mass. Arizona Washington Iowa Kazakhstan Arkansas

Work for multi-set?

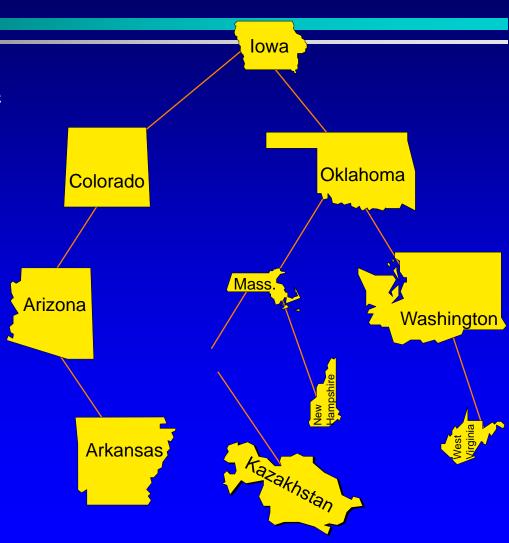
☐ If necessary, do some rearranging.

...copy that smallest item onto the item that we're removing...



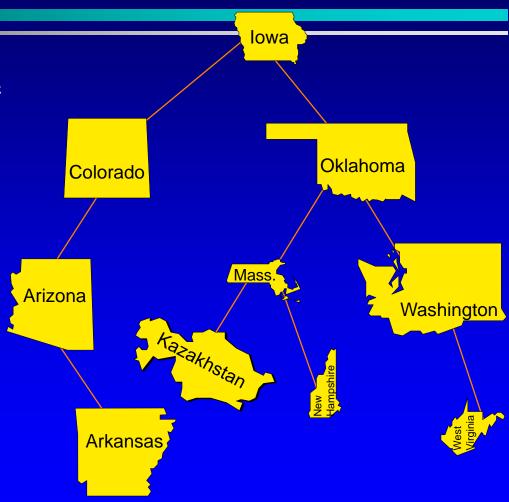
☐ If necessary, do some rearranging.

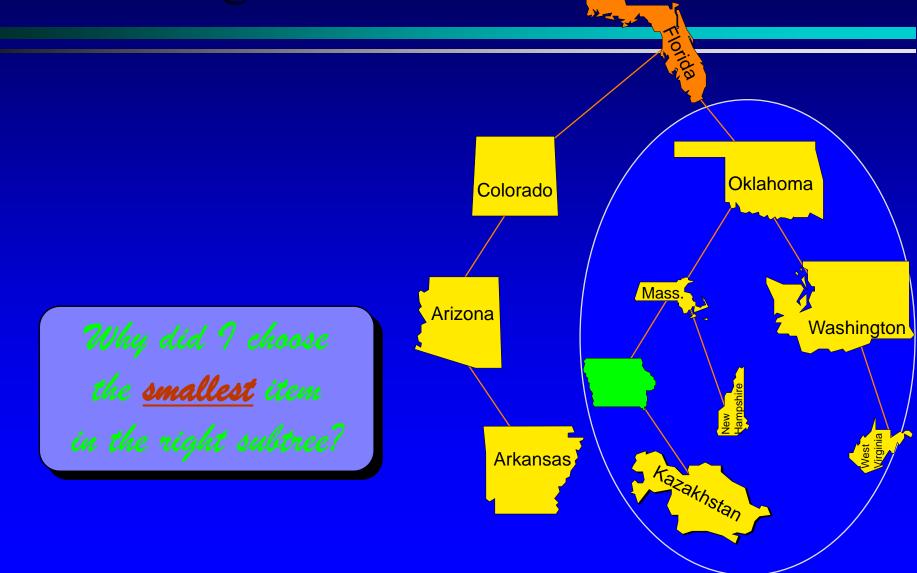
... and then remove the extra copy of the item we copied...

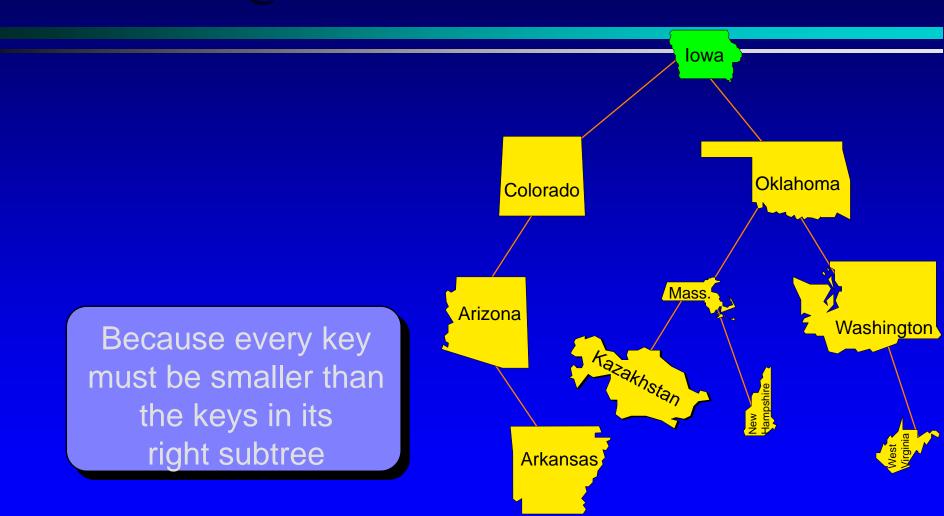


☐ If necessary, do some rearranging.

... and reconnect the tree







Removing an Item with a Given Key

- ☐ Find the item.
- ☐ If the item has a right child, rearrange the tree:
 - ☐ Find smallest item in the right subtree
 - □ Copy that smallest item onto the one that you want to remove
 - ☐ Remove the extra copy of the smallest item (making sure that you keep the tree connected)
 - else just remove the item.

Summary

- □ Binary search trees are a good implementation of data types such as sets, bags, and dictionaries.
- □ Searching for an item is generally quick since you move from the root to the item, without looking at many other items.
- Adding and deleting items is also quick.
- But as you'll see later, it is possible for the quickness to fail in some cases -- can you see why?

Assignment

- □ Read Section 10.5
- □ Assignment Bag class with a BST
 - Memeber functions
 - void insert(const Item& entry);
 - size_type count (const Item& target);
 - □ Non-member functions
 - viod bst_remove_all(binary_tree_node<Item>*& root const Item& target);
 - void bst_remove_max(binary_tree_node<Item>*& root, Item& removed);

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